

SPR 4106 Syntax and semantics in formal terms

Chapter 6 “Quantifiers”: 5 Essentials

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Quantificational DPs denote sets of sets of individuals

- For example (ignoring singular/plural difference):

$$\begin{aligned} \llbracket \text{most students} \rrbracket_s &= \\ \{ \llbracket \text{women} \rrbracket_s, \llbracket \text{diligent} \rrbracket_s, \llbracket \text{procrastinate} \rrbracket_s, \dots \} \end{aligned}$$

Quantificational determiners denote relations between such sets

- For example: $\llbracket \text{most} \rrbracket_s = \{ \langle X, Y \rangle : \text{most } X \text{ are } Y \}$
- Composition Principle: **generalized plugging in:**

$$\llbracket a + b \rrbracket_s = \llbracket a \rrbracket_s * \llbracket b \rrbracket_s = \{ z : \langle \llbracket b \rrbracket_s, z \rangle \in \llbracket a \rrbracket_s \}$$

if a denotes a relation between two individuals **or two sets**

Referential terms are also DPs – can they also denote sets of sets?

Yes they can: if their extensions are ‘lifted’ –
from the individuals themselves to the sets of sets containing them

Beside the name extension (i) we can define the DP extension (ii):

$$(i) \llbracket \text{Emilia} \rrbracket_s = e$$

$$(ii) \llbracket \text{Emilia} \rrbracket_s^{DP} = \text{LIFT}(\llbracket \text{Emilia} \rrbracket_s) = \{ X : e \in X \}$$

The result is the same:

$$\llbracket \text{Emilia rules} \rrbracket_s = 1 \text{ if and only if}$$

$$\llbracket \text{rules} \rrbracket_s \in \{ X : e \in X \} \text{ iff}$$

$$e \in \llbracket \text{rules} \rrbracket_s$$

Generalizing and simplifying the composition principles:

$$\llbracket a + b \rrbracket_s =$$

$$\begin{cases} \llbracket b \rrbracket_s(\llbracket a \rrbracket_s) & \text{if } \llbracket a \rrbracket_s \text{ is type } x \text{ and } \llbracket b \rrbracket_s \text{ is type } \langle x, y \rangle \\ \llbracket a \rrbracket_s \cap \llbracket b \rrbracket_s & \text{if } \llbracket a \rrbracket_s \text{ and } \llbracket b \rrbracket_s \text{ are type } \langle x, t \rangle \end{cases}$$

for any x and y , in particular, $x = e$ or $\langle e, t \rangle$ and $y = t$

- Just two principles: function application and intersection
- Equivalence between **set** and its **characteristic function**
- Function application **subsumes plugging in**: equivalence
 - between extension as a set of pairs $\langle u, v \rangle$ and
 - extension as a function from u to a set of v

Quantificational DPs in Object Position: Quantifier Raising (QR)

A **transitive verb** denotes a relation between individuals (or equivalently a function from individuals to sets of individuals)

A **quantifying DP** denotes a set of sets of individuals

Composing a type $\langle e, \langle e, t \rangle \rangle$ extension and a type $\langle \langle e, t \rangle, t \rangle$ extension would require a separate composition principle after all,

unless ...

- we use Quantifier Raising (QR), a covert movement operation
- $\llbracket \llbracket \text{both shoes} \rrbracket_i ; I \text{ like } t_i \rrbracket_s = \{x : I \text{ like}_s x\} \in \llbracket \llbracket \text{both shoes} \rrbracket_s$

Copula Verb *be* and Indefinite Article *a*

Indefinites are usually assumed to denote existential quantifiers:

$$\llbracket \text{a mitten} \rrbracket_s = \{ X : X \cap \llbracket \text{mitten} \rrbracket_s \neq \emptyset \}$$

And the copula *be* can be assumed to denote identity:

$$\llbracket \text{the winner is Estonia} \rrbracket_s = 1 \text{ iff } \llbracket \text{the winner} \rrbracket_s \in \{ x : x = e \}$$

Cases like (1) **can** be derived by combining these two assumptions:

(1) Scrooge is a miser.

But the right result can also be obtained by ignoring both *is* and *a*:

$$\llbracket \text{Scrooge (is a) miser} \rrbracket_s = \llbracket \text{miser} \rrbracket_s(s)$$